A Taxonomy for Task Allocation Problems with Temporal and Ordering Constraints

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Abstract

Previous work on assigning tasks to robots has proposed extensive categorizations of allocation of tasks with and without constraints. The main contribution of this paper is a more specific categorization of problems that have both temporal and ordering constraints. We propose a novel taxonomy that builds on the existing taxonomy for multi-robot task allocation and organizes the current literature according to the temporal nature of the tasks. We summarize widely used models and methods from the task allocation literature and related areas, such as vehicle routing and scheduling problems, showing similarities and differences.

1. Introduction

What is multi-robot task allocation? Think of a shipping company that sells an item every hour; a robot at the warehouse could receive that order, fetch the item, pack it, and prepare it for pick-up by a postal service. What happens when the company sells 20 items every hour? What about 20 items every minute? What about 20 items a second? Amazon, a popular shopping website, sold 36.8 million items on an especially popular shopping day in 2013. With 426 items ordered per second that day, a single robot would be hard-pressed to keep up with the orders. If the warehouse used a large team of robots, each robot would have to plan an efficient route through the warehouse to fetch items for shipping without colliding with other robots, without taking items that another robot is handling, all while planning its route around fetching items that are out-of-stock but will be restored soon.

Allocation of tasks with constraints on when, where, and in what order they need to be done by groups of robots is an important class of problems with many real-life applications. Applications include warehouse automation, pickup and delivery, surveillance at regular intervals, space exploration, search and rescue, and much more.

The nature of the temporal constraints in this class of problems is very broad; for example, in search and rescue domains the tasks are discovered over time and have to be done as quickly as possible. In dynamic environments, robots might
end up arriving late to some tasks and even miss some. On the other hand, success in surveillance tasks requires not to arrive to tasks late. Additionally, tasks may need to be executed in a specific order, such as in urban disaster scenarios in which police must clear blockades from roads before firetrucks can find and put out fires. Other tasks may need to be done synchronously, as in surveillance where robots have to track multiple people at the same time. The different nuances of temporal and ordering constraints lead to different models and solutions.

Previous taxonomies, such as Gerkey and Matarić [2004]’s key taxonomy and more recently the iTax taxonomy [Korsah et al., 2013], devote only limited attention to tasks that have temporal and ordering constraints. Our work attempts to fill this gap. We address the following research questions:

- What are the predominant types of temporal constraints in multi-robot task allocation?
- What are the most commonly used optimization objectives? Are they predominantly temporal-based, distance-based, or multi-objective?
- What models and methods from related areas can be applied to this class of problems?
- Which questions or variants have been answered well, and which remain largely open in this class of problems?

Our main contribution is a novel taxonomy that extends Gerkey and Matarić [2004], Korsah et al. [2013] taxonomies according to the nature of the temporal and ordering constraints considered. Gerkey’s taxonomy is based on three main characteristics of robots, tasks, and time. It considers the following axes:

- **Single-task robots (ST) vs. multi-task robots (MT)**: ST robots can do at most one task at a time, while MT robots can work on multiple tasks simultaneously.
- **Single-robot tasks (SR) vs. multi-robot tasks (MR)**: SR tasks require exactly one robot in order to be completed, while multiple robots are needed to complete an MR task.
- **Instantaneous (IA) vs. time-extended (TA) assignments**: In IA, tasks are allocated as they arrive, while in TA, tasks are scheduled over a planning horizon (defined in Section 2.1).

Our taxonomy focuses on time-extended task allocation problems with temporal and ordering constraints. Consequently, we drop the IA and TA distinction and replace it with our classification axes. However, when appropriate, we highlight if a problem or solution is of IA nature.
1.1. Organization

We begin by defining the class of multi-robot task allocation problems with temporal and ordering constraints (MRTA/TOC) in Section 2. In Section 3 we relate this class of problems to problems in other areas, setting the ground for our exploration of the models and methods in those areas. In Section 4 we present commonly used temporal and ordering models. In Section 5 we review the most common optimization objectives considered in the literature. Our taxonomy is introduced in Section 6. Task execution and the dynamics therein are discussed in Section 7. Solutions are introduced in Section 8. We discuss open issues, future directions, and final thoughts in Section 9.

Next, we formally define the task allocation problem with temporal and ordering constraints (MRTA/TOC), and summarize the terminology we use.

2. MRTA/TOC: Multi-robot Task Allocation with Temporal and Ordering Constraints

2.1. Terminology and Abbreviations

We define the terminology we use informally as follows:

- A robot is an autonomous agent responsible for performing some actions. Alternative names for robots are physical agents, unmanned vehicles, and rovers. Robots in MRTA/TOC are often modeled as holonomic and point robots, since the focus is not in low level control of robot motion.

- A team is a set of robots that work together. A team is called a coalition when it is formed to do some tasks and disbanded after that [Parker and Tang, 2006].

- A task is an action to be performed, also referred to as a work unit, activity, waypoint, or customer request. In some scheduling literature tasks are divided into jobs [Davis and Burns, 2011], while in other cases jobs consist of tasks [Balas et al., 2008].

- A time window is a time interval composed of the earliest time a task can start, and the latest time it can end. If an earliest time is not given, the latest time is referred to as a deadline constraint.

- A schedule is a timetable in which each task has a specific time to start, end, or both. In some cases each robot has its own individual schedule, while in others all robots share a single schedule.

- The scheduling horizon is the time period for which schedules are created. Alternatively, it is the end time, after which robots are not allowed to start or end tasks.

- The planning horizon is the time period over which plans are created.
The makespan is the time difference between the end of the last task and the start of the first task.

A route is a sequence of locations to visit. Routes and schedules are often used interchangeably, but schedules always concern time, while routes concern physical locations.

A task release refers to a task becoming available for execution. Task release can be deterministic if the release time is known upfront, dynamic if the release time is stochastic, or sporadic if it is governed by unknown probabilities; task release is also called periodic when the same task is released at regular intervals.

We use the following acronyms:

- MRTA/TOC for Multi-Robot Task Allocation with Temporal and Ordering Constraints.
- MIP for Mixed Integer Programming and MILP if the objective function and constraints are linear.
- TOPTW for Team Orienteering Problem (TOP) with Time Windows.
- VRPTW for Vehicle Routing Problem (VRP) with Time Windows.
- JSP for job-shop scheduling problems.

### 2.2. Problem Formulation

All the problems herein considered allocate tasks with either temporal or ordering constraints, or both. In MRTA/TOC we assume there is a finite set of robots and a set of tasks. A robot may have a location, velocity, route, and/or schedule. A task is defined by a subset of the following parameters: location, expected duration, cost, demand, reward, earliest start, and latest finish time.

Ordering constraints express a dependency between tasks, and are usually encoded as directed acyclic graphs. Each node in the graph represents a task, and each edge indicates precedence or simultaneity in the order of execution of the tasks.

The objective is to optimize some function of the cost (or reward) for doing the tasks for all the robots. Cost can be a time measure (e.g. makespan), or a spatial measure (e.g distance traveled). Commonly used optimization functions are more thoroughly described later in Section 5.

### 3. Connections with Other Problems

Multi-robot task allocation (MRTA) started in earnest in the 90’s, when researchers started pulling together teams of robots to accomplish multiple tasks. MRTA draws from a variety of areas in mathematics and operations research.
as well as computer science and robotics, including assignment problems, distributed computing, distributed AI, and scheduling.

The search for robust approaches to MRTA and related problems focused on how the robots perform in complex environments, leading researchers to add features like uncertainty with probabilistic and stochastic models, time windows for tasks, and spatial constraints. Solutions take different approaches, such as auctions, market-based planning, Markov Decision Processes, decentralized scheduling algorithms, and distributed constraint optimization.

In this paper we cover a subset of MRTA problems, which we call MRTA/TOC, to highlight the importance of temporal and ordering constraints among tasks and to shed light on how the inclusion of temporal and ordering constraints increases the complexity of task allocation.

Similar types of problems include the vehicle routing problem [Dantzig and Ramser, 1959], the job shop scheduling problem, and the team orienteering problem. Overall, multi-robot task allocation diverges from each of these problems on key points, including assumptions on the number of robots, robot and task homogeneity, environment dynamics caused by failures or interference from other robots, and communication restrictions.

We are now prepared to discuss the relationship between MRTA/TOC problems and vehicle routing problems with time windows (VRPTW), team orienteering problem with time windows (TOPTW), and job-shop scheduling problems (JSP).

**MRTA/TOC vs. VRPTW:** like the MRTA/TOC problem, the vehicle routing problem with time windows (VRPTW) [Kolen et al., 1987, Solomon and Desrochers, 1988, Desrochers et al., 1988, Toth and Vigo, 2002] studies problems which require solving allocation, routing, and scheduling subproblems simultaneously. Vehicles and robots are often treated as points in space, ignoring kinematic constraints, but kinematic [e.g. Cheng et al., 2008, for unmanned aerial vehicles] and sometimes dynamic [Pecora and Cirillo, 2012, for ground vehicles] constraints can be considered.

The solutions to several variants of VRPTW – such as multi-depot [Kang et al., 2005, Policek et al., 2004], dynamic and stochastic [Taş et al., 2013, Pavone et al., 2011, Laporte et al., 1992], and precedence and synchronization constrained [Korsah et al., 2012, Bredström and Rönnqvist, 2008] – have been extended to MRTA/TOC settings. An example of VRP similarities is the online pickup and delivery problem with transfers, where a team of vehicles has to pick up a set of items at a location and deliver them to another location [Coltin and Veloso, 2014a]. This problem is a generalization of the pickup and delivery problem which is well studied in operations research. However, the proposed solution is a typical MRTA approach. The authors combine a centralized temporal planner, which creates initial schedules, with auctions, which are used to repair the plans when delays or failures occur. In the same vein, Korsah et al. [2012] studied a MRTA problem that can be framed as a vehicle routing problem with temporal, precedence and synchronization constraints. The authors offer a MILP-based model and an optimal Branch-and-Price solution.

Despite their similarities, these problems differ in some ways. First, VRPTW
assumes an infinite number of vehicles is always available, with a few exceptions [e.g. Lau et al., 2003]). This assumption is not practical in robotic systems where the number of robots is usually fixed and can even decrease due to failures. VRPTW problems usually assume that all vehicles start from the same depot and return to the depot after work. In MRTA/TOC problems, robots may start at different locations and do not need to return to their initial locations. VRPTW problems mostly assume homogeneous vehicles with respect to their capabilities and capacities [for exceptions, see Bettinelli et al., 2011, Dondo and Cerdá, 2007]. In MRTA/TOC, however, robots are not necessarily homogeneous and their capacities and types can differ [Ponda et al., 2010, Schneider et al., 2005, Xu et al., 2005]. Lastly, MRTA/TOC problems, unlike VRPTW problems, usually use communication, often with constraints. In [Mercker et al., 2010] the communication graph is unknown (hence the algorithm does not always converge), while in [Ponda et al., 2012a] the communication graph is maintained by using specialized robots or robots not working on a task to act as communication relays. While in the previous two works convergence is guaranteed only for complete communication graphs, Jackson et al. [2013] and Smith and Bullo [2007] proposed distributed algorithms that converge with only local communication.

MRTA/TOC vs. TOPTW: in TOPTW, an origin and destination pair is given, and the goal is to search for control points to visit between the origin and destination such that the profit (or score function) is maximized while respecting all constraints. Each control point is associated with a profit (or score), and each edge connecting control points is weighted by the cost of moving between the control points [Labadie et al., 2012]. Control points are equivalent to tasks for robots in MRTA/TOC.

When TOPTW considers the origin and destination pairs to be the same point, then we have sub-tours similar to those for VRPTW problems, which can be described as vehicle routing problems with profit [Archetti et al., 2014]. One application of TOPTW problems, dial-a-ride, has gained some popularity in MRTA/TOC [e.g. Cottin and Veloso, 2014a, Rubinstein et al., 2012, Bouros et al., 2011]. In dial-a-ride, the problems are over-constrained [Carrabs F., 2007, Cordeau and Laporte, 2007], which means that not all the tasks can be performed, and thus the goal is to find the subset of tasks that maximizes the total profit [Rubinstein et al., 2012].

MRTA/TOC vs. JSP: job-shop scheduling problems are concerned with allocating groups of activities, called jobs, to a set of machines with the goal of minimizing the cost of completing the jobs, alone or in combination with other objectives [Allahverdi et al., 2008, Graham et al., 1979]. The problem can be decomposed into sequencing the activities and assigning start and end times to them (scheduling), which are solved simultaneously. Certain MRTA/TOC problems can be modeled as job-shop scheduling with setup times, deadlines and precedence constraints [Cesta et al., 2000, Balas et al., 2008, Oddi et al., 2011]; these problems include [Nunes and Gini, 2015, Gombolay et al., 2013, Dahl et al., 2009], although [Nunes and Gini, 2015] does not consider precedence constraints. In order to model MRTA/TOC problems as job-shop scheduling
problems, tasks are treated as jobs and robots as machines. We can map simple tasks to a job with only one activity, and complex tasks with subtasks to a job with multiple activities.

The mathematical models for job-shop scheduling do not apply directly to MRTA/TOC problems, because job-shop scheduling does not account for travel time. When setup times are used in job-shop scheduling, the setup time typically depends on the machine and not on the time needed for the job to reach the machine. The equivalent of travel time would be to use setup times that depend on the specific job [Korsah et al., 2013].

Modeling MRTA/TOC problems as job-shop scheduling problems is most useful when models and methods developed for scheduling [Cesta and Oddi, 1996, Cesta et al., 1999, Lee et al., 2009, Shah et al., 2009] are combined with MRTA solution techniques. In [Gombolay et al., 2013] a centralized approach is proposed in which a central temporal network is used and integrated with a MILP-based planner yielding near optimal schedules. In the decentralized approach of Barbulescu et al. [2010] each robot forms its own simple temporal network [Dechter et al., 1991], encoding both temporal constraints and precedence constraints in the network. To enforce precedence constraints a robot has to know which other robots depend on its schedule, so a high communication overhead is required to keep all robots up-to-date. The distributed approach in [Nunes and Gini, 2015] cuts down on communication costs by having each robot keep its own independent local temporal network and uses sequential single-item auction for allocation.

Having outlined the differences and similarities between MRTA/TOC and related problems, we now turn our attention to temporal and ordering constraints on MRTA problems.

4. Temporal Models and Task Ordering

Time models are outlined in Subsection 4.1 and Subsection 4.2, the nature of ordering constraints is presented in Subsection 4.3, while in Subsection 4.4 we discuss the nature of temporal constraints.

4.1. Relationships between time intervals

In general terms, time can be modeled using points or intervals [Allen, 1983]. An example time point is 10 am, while an interval is a continuous set of values bounded below and above by some time point, for example [10 am-12 pm]. When representing temporal constraints we may use either representation; however, the interval representation is much more common and is referred to as a time window. Time windows of tasks in general are allowed to overlap [e.g. Barbulescu et al., 2010, Gombolay et al., 2013, Heilporn et al., 2010, Koes et al., 2005, Nunes and Gini, 2015, Ponda et al., 2010, Luo, 2014].

The seminal paper by Allen [1983] proposed a set of relationships that hold between any two time intervals, as depicted in Fig. 1.

While the relationships originally were described between qualitative time intervals, they are also useful to describe the ordering between quantitative time
intervals. The relationships can be used to model partial or complete ordering
constraints between tasks, – for example, task X should be done before, after,
or at the same time as task Y. The X before Y operator can be used to describe
precedence constraints between tasks, while the X equal Y operator describes
a simultaneity constraint between the intervals or time points of two tasks.

4.2. Simple Temporal Networks (STN)

Equally influential is Dechter’s approach [Dechter et al., 1991], which pro-
posed to represent temporal constraints with a graph, called a simple temporal
network (STN). An example is in Fig. 2.

Figure 2: A simple temporal network with three tasks 1, 2 and 3. The self-loops on tasks
indicate the absolute start and end times for the task. Task 1 is done first, and then there is
a choice of doing 2 and 3 or just 3. \( st \) and \( ft \) are the actual start and finish times for each
task, \( es \) and \( ls \) are the earliest and latest times tasks can start, similarly \( ef \) and \( lf \) are the
earliest and latest times tasks can finish, and \( du \) represent tasks’ durations. \( tc_{k,k'} \) is a time
cost defined as the sum of the travel and wait times, which constrain when the next task \( k' \)
can be started.

Nodes represent time point variables or time events, and weighted edges
represent inequality constraints between time points. To reduce computational
complexity, this model requires exactly one constraint between every pair of
time point variables. This allows a solution to the scheduling problem to be
computed in polynomial time using the Floyd-Warshall algorithm. In an STN, the relationship between time windows can be represented by establishing constraints between start and finish times of tasks. While there are more complex models, for instance [Stergiou and Koubarakis, 2000, Block et al., 2006], in general these are NP-hard and can be approximated by solving several simple temporal problems [Boerkoel and Durfee, 2013].

STNs are commonly used in MRTA problems [Nunes and Gini, 2015, Gombolay et al., 2013, Barbulescu et al., 2010] because constraint consistency can be efficiently verified in polynomial time [Planken et al., 2008, Xu and Choueiry, 2003, Dechter et al., 1991]. An important feature of STNs is that new time points and constraints can be dynamically added in polynomial time [Coles et al., 2009, Cesta and Oddi, 1996], which is beneficial in dynamic domains where new tasks can appear and disappear.

STNs have been successfully extended to multi-agent settings [Boerkoel and Durfee, 2012, Boerkoel and Planken, 2012, Hunsberger, 2002] and to scenarios with uncertainties. Vidal [1999] uses set bounded uncertainty to model duration uncertainty of temporal events in an STN, and introduces the STN with uncertainty (STNU). Tsamardinos [2002] and Fang et al. [2014] extend STNUs by modeling uncertainty as probabilities. The former attempts to minimize the risk of temporal inconsistencies occurring, and the latter attempts to bound the probability of not meeting a schedule, respectively.

4.3. Task ordering

Precedence and simultaneity constraints are common in MRTA/TOC problems [Korsah et al., 2010, Gombolay et al., 2013, Barbulescu et al., 2010]. Ordering constraints force MRTA/TOC solutions to follow the partial or complete ordering of the tasks. Any solution violating the ordering of tasks is considered illegal. If time windows are pairwise disjoint, except possibly for the endpoints [Melvin et al., 2007], and the robots move in 2D then the strict order of the tasks simplifies the solution and allows for special cases where a polynomial solution exists.

The introduction of temporal and ordering constraints increases the complexity of task allocation, because solutions might contain assignments of tasks that depend on each other to different robots, creating execution dependencies among robots. This is undesirable because exogenous events affecting one robot will also affect all the robots that depend on the affected robot. The complexity of ordering constraints is further discussed in Jones et al. [2011], where intra-path precedence constraints among robots have been shown to add complexity to time-extended coordination solutions, when more than one task is assigned to each agent in domains that have intra-path constraints, such a disaster response. Luo et al. [2011] present an alternative model that divides tasks into disjoint sets with precedence constraints between the sets, each task takes the same time, and each robot can do only one task from each set.

A rare case is when executing a task precludes the execution of another. This type of problem is typically addressed at the planning stage, enforcing precedence constraints between the tasks [e.g. Olawsky and Gini, 1990].
4.4. Hard vs. Soft Constraints

Temporal constraints can be characterized as hard or soft constraints. Hard temporal constraints require that no temporal constraint is violated [Borning et al., 1992]. MRTA/TOC and related areas frequently require rigid time windows for tasks such as surveillance, routing for blood supply, and order fulfillment by warehouse robots, so many works focus on hard temporal constraints.

Soft temporal constraints allow some temporal constraints to be violated or some tasks to be skipped entirely, as long as the robot incurs a penalty for doing so [Bistarelli et al., 2007, Domshlak et al., 2006, Gerevini and Long, 2005].

Common types of soft temporal constraints include:

1. agents can start tasks early and/or finish tasks late with some penalty (called soft constraints in real time system terminology);
2. comply with the deadlines with some probability [Zheng and Woodside, 2003];
3. skip a number of consecutive tasks entirely or skip some percentage of tasks entirely [Bernat et al., 2001] without penalty (called weakly hard constraints in the real time systems terminology);
4. finish a task late without reward, or skip without penalty (called firm tasks [Bernat et al., 2001]);
5. use a mix of positive and negative preferences as constraints [Bistarelli et al., 2007, Domshlak et al., 2003], usually found in constraint and logic programming research.

The penalty incurred may differ depending on which constraint was violated; for example, finishing tasks late may be penalized more severely than doing tasks early. All non-hard constraints are called soft. Soft here is equivalent to modeling constraints as preferences, which encompasses all the non-hard constraints. These preferences do not take temporal constraints into account, but provide a way of comparing plans when several agreeable solutions are available.

Having discussed temporal models and constraints, and the nature of ordering constraints, we switch focus to optimization objectives. Determining these objectives is another important aspect to consider when building models for MRTA/TOC problems.

5. Optimization Objectives

Applications of MRTA/TOC problems require the robots to achieve a given optimization objective. In the rest of this paper we will refer to \( f(\cdot) \) as a generic function representing one of these objectives. There can be a single or multiple objectives [Jozefowiez et al., 2008]. Depending on the deterministic or stochastic nature of the problem, objectives will either be over actual or expected values. Optimization objectives might require a quantity to be minimized, usually a cost [Nunes and Gini, 2015, Gombolay et al., 2013, Chopra and Egerstedt, 2012] or regret [Heap and Pagnucco, 2014, Wu and Jennings, 2014], or to be maximized,
usually a score [Mercker et al., 2010, Ponda et al., 2010] or a reward [Korsah et al., 2012, Melvin et al., 2007, Koes et al., 2005]. Single optimization objectives may be of spatial nature (e.g. minimize total distance traveled) or of temporal nature (e.g. minimize makespan).

Common optimization objectives for MRTA/TOC problems include:

- **MiniSUM**, i.e. minimize the sum of the robot path costs over all the robots [Lagoudakis et al., 2005]. Minimizing the distance traveled is common [e.g. Coltin and Veloso, 2014b, Chopra and Egerstedt, 2012, MacKenzie, 2003]) but some instead minimize a time measure over robot paths [e.g. Heap and Pagnucco, 2014, Barbulescu et al., 2010]).

- **MiniMAX**, i.e. minimize the maximum path cost of a robot over all the robots [Lagoudakis et al., 2005]. Instead of minimizing the maximum path cost, a similar objective function is to minimize the makespan, i.e. the difference between the start of the first and the end of the last task [Graham et al., 1979]. In [Nunes and Gini, 2015] the makespan is minimized in a decentralized manner while in [Gombolay et al., 2013] the makespan, along with other objectives, is minimized using a near-optimal centralized MILP-based planner.

- **MiniAVE**: i.e. minimize over all the tasks the average cost of the path for a robot from its initial location to the task location, assuming each task is visited by a single robot [Lagoudakis et al., 2005]. This is known as the Traveling Repairman Problem [Fakcharoenphol et al., 2007], where the objective is to minimize the wait time of the customers (or tasks) for a repairman (or robot).

- **Minimize lateness or tardiness**, which is the difference between the earliest start time of a task and the actual arrival time of the robot [Ponda et al., 2010, Rubinstein et al., 2012, Beck and Refalo, 2003]. A similar objective is to minimize the idle time of the robots [Hasgül et al., 2009].

- **Maximize the number of tasks completed** [Lau et al., 2003, Colorni and Righini, 2001] or minimize the number of tasks missed [Hasgül et al., 2009].

- **Minimize the number of robots used.** This is very common in vehicle routing problems, where there is an unlimited number of vehicles available [Luo and Schönfeld, 2007, Bräysy and Gendreau, 2005a, Desrochers et al., 1988].

- **Maximize profit**, measured as the difference between the reward of tasks and their respective costs [Korsah et al., 2012, Melvin et al., 2007], or as the team utility [Amador et al., 2014, Ponda et al., 2010, Koes et al., 2005].

While not extensively covered here, multi-objective problems are common [Jozefowiez et al., 2008], especially when objectives are combined through linear aggregation. For example, makespan and distance are minimized in [Ponda...
et al., 2010, Nunes and Gini, 2015], while [Gombolay et al., 2013] also minimizes workspace overlap.

The problems in the subcategories of our taxonomy can, in most cases, be formalized using MIP programs. In other cases, where tasks’ locations or durations, or travel times are probabilistic, stochastic models (e.g. Markov Decision Processes) are more commonly used. We give specific examples of different objectives within the subcategories, using \( f(\cdot) \) as a generic objective function. The constraints include coverage constraints that dictate the number of robots required to complete a task as well as the number of tasks a robot is allowed to complete at a time; ordering constraints, for example precedence and simultaneity constraints; and side constraints, such as resource constraints.

We summarize the notation we use in Table 1.

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<thead>
<tr>
<th>Agents</th>
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<tbody>
<tr>
<td>( A )</td>
<td>set of agents</td>
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<tr>
<td>( a )</td>
<td>agent in set ( A )</td>
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<tr>
<td>( q_a )</td>
<td>capacity of agent ( a )</td>
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<tr>
<td>Tasks</td>
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<td>( K )</td>
<td>set of tasks</td>
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<td>( k )</td>
<td>task in set ( K )</td>
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<td>( es_k )</td>
<td>earliest start time of task ( k )</td>
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<tr>
<td>( ls_k )</td>
<td>latest start time of task ( k )</td>
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<td>( ef_k )</td>
<td>earliest finish time of task ( k )</td>
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<td>( lf_k )</td>
<td>latest finish time of task ( k )</td>
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<td>actual start time of task ( k )</td>
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<td>( ft_k )</td>
<td>actual finish time of task ( k )</td>
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<tr>
<td>( du_k )</td>
<td>duration of task ( k )</td>
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<tr>
<td>Optimization</td>
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<tr>
<td>( f(\cdot) )</td>
<td>generic optimization function</td>
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<td>( x_k^a )</td>
<td>indicator assignment of task ( k ) to robot ( a )</td>
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<td>( o_{k,k'}^a )</td>
<td>indicator that agent ( a ) performs task ( k' ) directly after ( k )</td>
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<tr>
<td>( tt_{kk'} )</td>
<td>travel time between tasks ( k ) and ( k' )</td>
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<tr>
<td>( w_k^a )</td>
<td>workload for task ( k ) when performed by robot ( a )</td>
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<tr>
<td>( v_k^a )</td>
<td>indicator that robot ( a ) performs task ( k ) first</td>
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<tr>
<td>( z_k^a )</td>
<td>indicator that robot ( a ) performs task ( k ) last</td>
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<tr>
<td>( Y_a )</td>
<td>set of possible routes for agent ( a )</td>
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<tr>
<td>( x_y^a )</td>
<td>indicator assignment of route ( y ) to agent ( a )</td>
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<tr>
<td>( b_{y,k}^a )</td>
<td>indicator that task ( k ) is in route ( y ) of agent ( a )</td>
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<tr>
<td>( c_y^a )</td>
<td>cost to agent ( a ) of route ( y )</td>
<td></td>
<td></td>
</tr>
<tr>
<td>( u_k^a )</td>
<td>reward agent ( a ) collects for performing task ( k )</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Table 1: Notation used in the paper
6. Taxonomy

We are now ready to introduce our extensions to the taxonomy of Gerkey and Matarić [2004] and focus on time-extended assignments, in which robots build schedules for the tasks. We categorize the literature according to the temporal constraint types, while keeping the ST-MT and SR-MR axes to describe the robot capabilities and the task needs, respectively. We add the following new axis in our taxonomy:

- **Hard temporal constraints (HC) vs. Soft temporal constraints (SC).** Hard temporal constraints require that no temporal constraint for any task is violated. Soft temporal constraints allow some temporal constraints to be violated or some tasks to be rejected entirely with a penalty.

We now illustrate our taxonomy in terms of single- vs. multi-task robots (SR - MR), single- vs. multi-robot tasks (ST - MT), and hard vs. soft constraints (HC - SC). Most categories are further broken down as deterministic vs. stochastic problems. We begin with the least complex problem settings, in which single-task robots are allocated to single-robot tasks.


6.1.1. Deterministic allocations

Deterministic ST-SR-HC problems typically assume that there are more tasks than robots, and that all tasks are known in advance; they require time-extended assignments. These problems are comprised of three intertwined subproblems: (1) an assignment subproblem, to find the assignment of tasks to robots that optimizes the given objective function $f(\cdot)$; (2) a task sequencing subproblem, to find feasible orderings of tasks that result in optimal assignments, and (3) a scheduling subproblem, to assign times to tasks in a way that optimizes $f(\cdot)$.

Tasks have to be scheduled so that no temporal or assignment constraints are violated. Temporal constraints are violated when robots do tasks at times that are not consistent with the temporal constraints on the tasks. Assignment violations occur when two or more robots are assigned to the same task, or two or more tasks are scheduled to be done at the same time by the same robot.

In the mixed integer linear programming formulation in Fig. 3, $x_k^a$ is an indicator variable that takes the value 1 if robot $a$ is assigned task $k$ and 0 otherwise, $o_{kk'}^a$ is an indicator variable that takes the value 1 if robot $a$ performs task $k$ followed directly by task $k'$, and 0 otherwise. $q_a$ is the capacity of robot $a$, $st_k$ and $ft_k$ are respectively the actual start and finish times for task $k$, $tt_{kk'}$ is the travel time between tasks $k$ and $k'$, $w_k^a$ is the amount of work robot $a$ has to perform when assigned task $k$, $v_k^a$ is a binary variable that is 1 if task $k$ is the first task in robot $a$’s schedule and 0 otherwise, and $z_k^a$ is a binary variable that is 1 if task $k$ is the last task in robot $a$’s schedule and 0 otherwise. We assume all times are strictly positive.

Since robots start at their own initial location, we create an empty task for each robot $a$ at its initial location. The empty task starts at time 0, and has...
minimize or maximize  \( f(\cdot) \)

subject to

(a) \( \sum_{a \in A} x_k^a = 1 \) \quad \forall k \in K_a^+ 

(b) \( \sum_{k \in K_a^+} v_k^a = 1 \) \quad \forall a \in A 

(c) \( \sum_{k \in K_a^+} z_k^a = 1 \) \quad \forall a \in A 

(d) \( \sum_{k \in K_a^+} w_k^a x_k^a \leq q_a \) \quad \forall a \in A 

(e) \( \sum_{k \in K_a^+} o_{kk'}^a + v_k^a = x_k^{a'} \) \quad \forall a \in A, k', k \in K_a^+ 

(f) \( \sum_{k' \in K_a^+} o_{kk'}^a + z_k^a = x_k^{a'} \) \quad \forall a \in A, k, k' \in K_a^+ 

(g) \( es_k \leq st_k \leq ls_k \) \quad \forall k \in K_a^+ 

(h) \( ef_k \leq ft_k \leq lf_k \) \quad \forall k \in K_a^+ 

(i) \( ft_k - st_k \geq du_k \) \quad \forall k \in K_a^+ 

(j) \( ft_k + tt_{kk'} - M * (1 - o_{kk'}^a) \leq st_{k'} \) \quad \forall a \in A, k, k', k' \in K_a^+ 

(k) \( x_k^a \in \{0, 1\} \) \quad \forall a \in A, k \in K_a^+ 

(l) \( o_{kk'}^a \in \{0, 1\} \) \quad \forall a \in A, k, k', k' \in K_a^+ 

(m) \( v_k^a \in \{0, 1\} \) \quad \forall a \in A, k \in K_a^+ 

(n) \( z_k^a \in \{0, 1\} \) \quad \forall a \in A, k \in K_a^+ 

For instance, to minimize the makespan the optimization function would be

\[
\text{minimize } f(t_k) \quad \forall k \in K_a^+, \forall a \in A 
\]

For ST-SR-HC problems the coverage constraints in Fig. 3 enforce that (a) each task gets at most one robot, each robot has a first (b) and last (c) task and that (d) each robot does as many tasks as its capacity allows. Capacity here means maximum workload a robot is allowed to perform.

The sequencing constraints require that (e) every task \( k \) assigned to robot \( a \) except the first has a predecessor, and that (f) every task except the last has a successor. Temporal constraints (g)–(j) are constraints on the service times of tasks. Constraint (j) ensures that the interval between two consecutive tasks is large enough for the robot to travel to it. The constraint includes a sufficiently large constant \( M \) to make the formulation a mixed-integer linear program. Constraints (k)–(n) bound the values for the indicator variables.
Ordering constraints in ST-SR-HC problems, such as synchronization constraints, do not differ significantly from other types of temporal constraints, thus we will only show one example of such constraints here. Eq. 2 and Eq. 3 are an example of the formulation of these constraints. Let $k, k' \in P$ where $P$ is a set of task pairs with ordering constraints, and $P^{\text{sync}} \subseteq P$ is the subset of tasks that have to be performed simultaneously, i.e. start at the same time. Eq. 2 states that regardless of which robot(s) is assigned to tasks $k$ and $k'$, task $k'$ should start $\epsilon$ time units after the finishing time of task $k$. If $\epsilon > 0$, $k, k' \in P$ (Eq. 2), and if $\epsilon = 0$ then $k, k' \in P^{\text{sync}}$ (Eq. 3).

$$\sum_{a \in A} s_{k'} x_{a}^{k} - \sum_{a \in A} f_{k} x_{a}^{k} > \epsilon + M(1 - o_{kk'}^{a}) \quad \forall a \in A, k, k' \in P, \epsilon > 0$$ (2)

$$\sum_{a \in A} s_{k'} x_{a}^{k'} - \sum_{a \in A} s_{k} x_{a}^{k} = 0 \quad k, k' \in P^{\text{sync}}$$ (3)

Advances in MILP formulations for VRPTW [Barnhart et al., 1998, Feillet, 2010] and more recently for MRTA problems [Korsah et al., 2012] have proposed set covering and set partitioning-based formulations. These formulations assign routes, instead of tasks, to robots. The allocation problem is decomposed into what is known as the master problem, and a pricing subproblem. One of the gains of such formulations is that the master problem can be restricted to evaluating subsets of tasks at a time, instead of the entire set of tasks. Pricing subproblems solve temporally constrained shortest path problems rooted at robot locations, in which routes can be computed via heuristic methods, such as D* lite as in [Korsah et al., 2012]. Such formulations benefit from the insight that for very large problems many routes are not part of any optimal solution. Thus, selectively incrementing candidate routes decreases computational and memory costs. Feillet [2010] provides a technically rigorous tutorial of such formulations and their advantages for VRPTW problems.

As defined in Table 1, let $Y_{a}$ be a set of routes for robot $a$ computed using the shortest path algorithm with resource constraints; $\hat{x}_{y}^{a}$ is an indicator variable that assumes a value of 1 if robot $a$ is assigned route $y \in Y_{a}$ and 0 otherwise; $C_{y}^{a}$ is the expected cost robot $a$ incurs for performing route $y$; finally, $b_{yk}^{a}$ is an indicator variable that is 1 if task $k$ is performed in route $y \in Y_{a}$ belonging to robot $a$ and 0 otherwise. An example of a set partitioning formulation of ST-SR-HC problems is shown in Fig. 4.

One of the main limitations of deterministic ST-SR-HC problems is their applicability. These problems do not take uncertainty and partial knowledge into account, and these are important properties of many robotics problems. Next, we discuss ST-SR-HC problems that handle uncertainty in planning by modeling uncertainty as stochastic processes.

### 6.1.2. Stochastic allocations

In stochastic ST-SR-HC problems, it is assumed that a model of uncertainty is available. Stochastic ST-SR-HC, like other stochastic problems in our taxonomy, are usually modeled as pure or mixed stochastic integer programs, or as
minimize \( \sum_{a \in A} \sum_{y \in Y_a} C^{a}_{y} \hat{x}^{a}_{y} \)
subject to
(a) \( \sum_{a \in A} \sum_{y \in Y_a} \hat{x}^{a}_{y} \leq 1 \) \( \forall a \in A \) 
Every robot gets only 1 route.
(b) \( \sum_{a \in A} \sum_{y \in Y_a} \hat{x}^{a}_{y} b^{a}_{yk} = 1 \) \( \forall k \in K \) 
Each task is on 1 route.
(c) \( \hat{x}^{a}_{y} \in \{0, 1\} \) \( \forall a \in A, y \in Y_a \) 
Indicator: route to robot.
(d) \( b^{a}_{yk} \in \{0, 1\} \) \( \forall a \in A, y \in Y_a, k \in K \) 
Indicator: task to route.

Figure 4: Set partitioning formulation example for MRTA/TOC problems.

Markov Decision Processes (MDPs) [Gendreau et al., 1996]. When modeled as
stochastic integer programs [Ponda et al., 2012b] they assume the form in Eq. 4
with the constraints shown in Fig. 3 or stochastic constraints [Shen et al., 2009].
In Eq. 4 the objective function is the expected reward, \( \theta \in \Theta \) is the uncertainty
model that is available to the robots, and \( u^{a}_{k} \) is the reward that agent \( a \) gets for
doing task \( k \).

\[
\text{maximize } E_{\theta}(\sum_{a \in A} \sum_{k \in K} u^{a}_{k} x^{a}_{k}) \tag{4}
\]

Examples of uncertainty models include probability distributions for task ar-
vival, robot travel time, task availability, and more [Miao et al., 1991]. Stochastic
formulations other than MDPs are not, to the best of our knowledge, widely
used in the MRTA literature, so we turn to the dynamic and stochastic VRPTW
literature for examples of such models. For instance, [Bopardikar et al., 2014]
studied a dynamic VRP problem in which demands (or tasks) with deterministic
time constraints arrive randomly. A Poisson process generates the time when a
task appears, while a uniform distribution is used for the demand location. The
main goal of the work is to maximize the fraction of demand met.

Similarly, in [Pavone et al., 2009] demand is stochastic; however, they also
consider time window constraints. They study stochastic and dynamic VRPTW
problems, with the objective of minimizing the number of utilized vehicles and
maximizing the demand satisfied. Both [Bopardikar et al., 2014] and [Pavone
et al., 2009] analyze a different number of requirements, such as bounds on the
number of vehicles used and maximum number of tasks that can be missed.
In both, temporal constraints cannot be violated. However, in order to prove
properties about their solutions, some strong assumptions are made, such as all
time windows have the same length [Pavone et al., 2009].

An alternative way of modeling uncertainty uses MDPs. In [Dean et al.,
1993, Beynier and Mouaddib, 2007] MDP states are locations in a map with
obstacles, tasks and robots. In [Beynier and Mouaddib, 2007] a state is a triplet
representing the previously visited state, the amount of resources left, and the
time window. The goal is to search for policies that maximize a value function for
the augmented states. Dolgov et al. [2007] poses the problem as a combinatorial
resource scheduling problem with uncertainty, which can be easily extended to
include locations, forming an MRTA problem.

Uncertainty models for ST-SR-HC problems are, to the best of our knowledge, rarely explored in the vast MRTA literature, although stochastic planning could lead to practical gains in terms of finding sound and robust allocation policies for robots. Instead, it is far more common to find papers that address stochasticity by model-free methods, such as reinforcement learning or that deal with uncertainty by simply replanning during task execution.


6.2.1. Deterministic allocations

Deterministic ST-SR-SC problems and deterministic ST-SR-HC problems are very similar, except that they differ in the hardness of the time constraints. Classic problems include vehicle routing and scheduling problems, task sequence, and soft constraint scheduling subproblems. We will also look at problems and solutions in the real time systems and artificial intelligence literature to assist in modeling temporal and ordering constraints in MRTA.

In soft time window constraints for vehicle routing, the goal is to find the best agent-task assignments that minimize the cost function \( f(\cdot) \) of servicing some number of clients. The total cost of assigning a set of agents or vehicles and departure times is equal to the fixed cost of operating the agents, plus the cost of operating the agents on the specific routes, plus the penalty cost for arriving early or late to the clients on the routes [Taş et al., 2013, Hsu et al., 2007, Ando and Taniguchi, 2006, Taillard et al., 1997]. Penalty costs for arriving early may be different than for arriving late, and these may vary by domain.

Frequent assumptions are that agents may operate more than one route a day, each client must be serviced exactly once and by one agent (see Fig. 3), clients have time windows in which to be serviced for a specific amount of time, and agents have capacity constraints.

Though some soft constraint types are not widely used in MRTA, such as the weakly hard constraints in real time systems, they transfer quite sensibly into MRTA problems. In a weakly hard system with periodic task release, the distribution of met and missed deadlines in a time period is precisely bounded [Bernat et al., 2001]. Other approaches to skipping some deadlines for periodic tasks include degradation policies in overloaded systems [Beccari et al., 1999] or exploiting skips to improve response time for aperiodic tasks [Caccamo and Buttazzo, 1997]. MRTA problems most frequently have sets of tasks that must be assigned and completed once; the tasks may be similar or near each other spatially or temporally, but in general, there are no repeatedly released tasks.

Real world robot-task assignment problems, however, might demand periodic tasks. A Mars rover, for example, has a regularly scheduled self-maintenance period, as well as periodic deadlines to finish uploading data or downloading instructions. These deadlines are usually hard deadlines, so the robot can shut down overnight and clear memory caches; other regularly scheduled robotic activities are not so sensitive to the time of execution.

A periodic task’s worst case met and missed deadlines can be considered by the number of random or consecutive missed or met deadlines [Bernat et al.,
There are four constraints to consider: making any $n$ in $m$ deadlines, making $n$ in a row in $m$ deadlines, missing any $n$ in $m$ deadlines, and missing $n$ in a row in $m$ deadlines. In this way, any regularly scheduled sequence of tasks can allow some missed deadlines without penalty, while still allowing the agent responsible for those tasks to schedule and make most of its deadlines. Agricultural drones, for example, may have regularly scheduled sampling, such as fertilization, weed picking, or soil testing responsibilities that allow to skip a few deadlines.

6.2.2. Stochastic formulations

Stochastic versions of ST-SR-SC problems have an uncertainty model available, as in Section 6.1; now, however, we use soft windows and allow agents to gain value even when performing tasks outside their original time window. Our objective is still to minimize the cost function given earlier in the deterministic formulation, with the inclusion of some probability model; often these are probabilities of travel delay between tasks and therefore travel times. The cost function now includes the cost of using vehicles and the cost of arriving to a task outside the proper time window.

Work from [Ta¸s et al., 2013, Ando and Taniguchi, 2006] distinguishes between fixed costs of operating a vehicle and service costs, which vary depending on the route and are impacted by the distribution of travel times. Work in [Ta¸s et al., 2013] models travel time delays with several distribution types, which change the variable service cost of operating a vehicle. Besides changing the cost of using a specific vehicle, stochastic formulations impact the arrival times and the cost of doing a task early or late. If we use soft time windows, we can vary the cost of arriving early (e.g. early arrival is a small penalty) or of arriving late (e.g., late arrival is a large penalty). The travel time probability directly impacts whether the agent arrives early or late, which is why we frequently see stochastic formulations in soft time windows but no other kinds of preferred constraints.


6.3.1. Deterministic allocations

In ST-MR-HC allocation problems, agents are scheduled to work simultaneously on tasks as coalitions. Coalition-based task allocation occurs when tasks cannot be executed by a single agent, or when task execution is more efficient when done by more agents [Vig and Adams, 2006, Shehory and Kraus, 1998]. In disaster rescue, for instance, fire fighters working in coalitions may extinguish the same number of fires earlier than if these rescuers had to work individually on each fire [Parker et al., 2015]. Moreover, in scenarios where the number of agents is limited, coalition-based allocations may enable a higher task completion rate [Ramchurn et al., 2010b].

Coalition formation, in general, requires dealing with two subproblems: coalition value computation and coalition structure generation [Sandholm et al., 1999]. The former is concerned with computing the expected utilities (or costs)
of forming all possible coalitions, whereas the latter is concerned with partitioning the set of agents into exhaustive and disjoint groups that maximize the total utility. In MRTA, the coalition value is typically a combination of the utility gained and the coordination cost necessary to perform a task. Coalition size may be restricted by the physical constraints which limit the number of agents that can simultaneously work on the same task.

Let $2^A$ be the set of agent coalitions that may be formed with the agents in $A$ (i.e., all subsets of $A$) and $x^c_k$ be an indicator variable that takes the value of 1 if coalition $c \in 2^A$ is assigned to task $k$ and 0 otherwise. For simplicity’s sake, we assume that all agents start their tours from an initial node 0 and finish at node $m + 1$. Let $o^a_{kk'} = 1$ when agent $a$ visits task $k'$ directly after $k$. $o^a_{0k} = 1$ denotes the fact that $a$ visits $k$ at the very beginning of the route, and 0 otherwise. Similarly, $o^a_{k(m+1)} = 1$ when agent $a$ visits task $k$ at the end of its route, and 0 otherwise. ST-MR-HC allocation problems can generally be formalized by the MILP in Fig. 5.

minimize or maximize $f(.)$

subject to

(a) $\sum_{c \in 2^A} x^c_k \leq 1$ \hspace{1cm} $\forall k \in K$
(b) $\sum_{a \in c} x^a_k = |c| x^c_k$ \hspace{1cm} $\forall c \in 2^A, k \in K$
(c) $\sum_{k \in K} o^a_{0k} = 1$ \hspace{1cm} $\forall a \in A$
(d) $\sum_{k \in K} o^a_{k(m+1)} = 1$ \hspace{1cm} $\forall a \in A$
(e) $\sum_{k \in K, k \neq k'} o^a_{kk'} - \sum_{k'' \in K, k' \neq k''} o^a_{k''k'} = 0$ \hspace{1cm} $\forall a \in A, k, k' \in K$
(f) $s_{tk} + du_k + tt_{kk'} - M \ast (1 - o^a_{kk'}) \leq s_{tk'}$ \hspace{1cm} $\forall a \in A, k, k' \in K$
(g) $e_{sk} \leq s_{tk} \leq l_{sk}$ \hspace{1cm} $\forall k \in K$
(h) $e_{fk} \leq f_{tk} \leq l_{fk}$ \hspace{1cm} $\forall k \in K$
(i) $f_{tk} - s_{tk} \geq du_k$ \hspace{1cm} $\forall k \in K$
(j) $x^a_k \in \{0, 1\}$ \hspace{1cm} $\forall a \in A, k \in K$
(k) $x^a_k \in \{0, 1\}$ \hspace{1cm} $\forall c \in 2^A, k \in K$
(l) $o^a_{kk'} \in \{0, 1\}$ \hspace{1cm} $\forall a \in A, k \in K, k' \in K$

Figure 5: Standard mixed integer formulation of the task allocation problem with single-task robots, multiple-robot tasks, and hard temporal constraints

Constraint (a) guarantees allocations of no more than one coalition per task. As the problem maybe over-constrained, not all tasks may be allocated. Constraint (b) guarantees if a coalition is assigned to a task then all the agents in the coalition are assigned to that task too. Constraint (c) guarantees that all the agents start from the initial location, and constraint (d) ensures that they finish their routes at the final location. Constraint (e) guarantees the connectivity of the routes, so that a robot reaches all its assigned tasks in sequence. Constraints (f)–(i) ensure that the visit time-line is feasible and the time windows are respected (as in MILP for ST-SR-HC in Fig. 3). An extra waiting time may be imposed after the task’s earliest start time to form the coalition. When coalition work affects the task execution efficiency, the task duration ($du_i$)
should be computed accordingly [Ramchurn et al., 2010b]. Constraints (j)–(l)
bound the values for the indicator variables.

An alternative model for ST-MR-HC problems is the set partitioning model.
When cast as a set partitioning problem, a set of coalitions $S = \{c_1, \ldots, c_{|S|}\}$
corresponds to a set partition of $A$ if and only if $\bigcup_{c_i \in S} c_i = A$ and the elements
of $S$ are pairwise disjoint (i.e., $\forall c_i, c_j \in S$ s.t. $i \neq j : c_i \cap c_j = \emptyset$). The
solution to the set partition problem is a partition of $A$ that maximizes the
utility $u : S \rightarrow \mathbb{R}^+$. The NP-hard nature of problems in this class require
approximate solutions for practical coalition-based MRTA problems.

Certain side constraints are very important in this class of MRTA/TOC
problems. Examples of such constraints include capability and resource con-
straints. Agents may have limited resources, especially in the case of small
robots. For instance, in disaster rescue scenarios, fire trucks may need a certain
amount of fuel to travel to a fire and a certain amount of water to extinguish it.
Tasks might require coalitions of agents with certain capabilities. For instance,
a fire might require a coalition of fire fighters, while police and ambulances can
collaborate to dig out and carry survivors to refuge centers [Kitano and Satoshi,
2001].

MRTA researchers have proposed coalition-based frameworks for heteroge-
neous robotics. Examples include the ASyMTRe [Parker and Tang, 2006] ar-
chitecture. ASyMTRe is a reasoning system for heterogeneous robots to form
coalitions to do tasks that require tight robot coordination. The architecture
uses a collection of schemas for perception and motor control, which are con-
nected at run time, enabling the robots to share information as needed to com-
plete the tasks. The architecture has been extended [Zhang and Parker, 2013a]
to ensure that only feasible coalitions are formed. Efficient scheduling heuristics
for coalitions are proposed in [Zhang and Parker, 2013b].

To the best of our knowledge, no literature addresses stochastic ST-MR
problems with either hard or soft constraints. We will now examine tasks with
soft constraints.


ST-MR-SC assumes that multiple agents can work simultaneously on the
same task and are allowed to violate some temporal constraints, as long as they
incur a penalty for the violation. ST-MR-SC can be likened to ST-MR-HC
problems. The difference is that in the ST-MR-SC's case the objective function
takes a temporal violation penalty into account (Eq. 5).

$$
\arg\max_{S \in \mathcal{P}(A)} \sum_{c \in S} \sum_{k \in K} x^k_c u(c, k) \pi(k, st_k) \tag{5}
$$

In Eq. 5 $S$ is a coalition structure, $u(c, k)$ is coalition $c$'s utility for performing
task $k$ and $\pi(k, st_k) \in [0, 1]$ is the utility decay coefficient function for task $k$.
This coefficient is set as $\pi(k, st_k) = 1$ when task $k$ is started and/or finished
within the time window (i.e., $\epsilon s_k \leq t \leq l s_k$ and $st_k + du_k \leq l f_k$). Early and/or
late task executions are penalized by setting $\pi(\cdot) \in [0, 1]$. In particular, we set $\pi(k, st_k) = 0$, $\forall k \in K$ when $st_k + du_k > ls_k$ [Koes et al., 2005, Amador et al., 2014]. Fig. 6 illustrates the difference between a soft deadline utility function (left) and a hard deadline utility function (right).

There is very little research on ST-MR-SC problems. Most of this research is motivated by application areas such as urban search and rescue [Koes et al., 2005, Scerri et al., 2005], or law enforcement where police officers are assigned to crime events in a city [Amador et al., 2014]. In [Scerri et al., 2005] an expected utility model is used to allocate interdependent tasks. Late task executions are penalized by subtracting the delay cost from the total utility. The work subdivides large tasks into smaller subtasks that are linked with simultaneous execution interdependency, and coalitions of agents execute the smaller subtasks. The coalition formation problem is simplified by fixing the coalition size and reducing the number of allowed coalitions.

In [Koes et al., 2005] the task utility decays over time from the beginning of the mission and becomes zero by the mission deadline. Likewise, in [Amador et al., 2014] the utility of tasks delayed beyond the soft deadline decays exponentially over time. The coalition value depends on the number of agents and is a function of the agents’ fitness in performing a task.


6.5.1. Deterministic allocations

The MT-SR problem with hard temporal constraints is no more common now than it was in [Gerkey and Mataric, 2004], but we can provide some additional context for multi-task robots regardless of hard or soft time constraints. Gerkey’s work likens the MT-SR problem to the ST-MR problem, using the same mathematical formulation for both problems but switching the role of tasks and agents in the formula. Multi-task robots do exist in real life, however, such as the mission-driven Mars rover Curiosity, so the multi-task robot prob-
lem merits deeper discussion before we consider it in terms of hard or soft time windows.

A multi-tasking robot can perform located tasks, such as grasping and manipulating objects, and non-located tasks, such as taking pictures of nearby objects. Furthermore, located tasks can be near to or far away from the robot; for example, an object two feet in front of the robot is close, but an object ten feet away is probably considered far from the robot. Nearby objects should be relatively easy to grasp, but farther-away objects will require larger or longer actuators and thus more complex kinematic calculations to properly manipulate them. Consider unmanned aerial vehicles; reconnaissance drones may track objects and take pictures (a relatively easy task) or may need to track objects on the ground and drop packages (a more difficult task that includes more intense object manipulation).

Another complexity for the multi-tasking robot arises in the solution methods used for task assignment, which will be addressed more thoroughly in Section 8. Are the tasks assigned in a centralized or de-centralized fashion? A centralized system can produce optimal assignments, but if the robot decides on its own which tasks to perform at once, it must schedule its own resources according to the time frame and load on its system.

Lastly, a multi-tasking robot can either preempt tasks or not; preemptable tasks require priority knowledge and may require task rescheduling, whereas a simpler system of non-preemptive tasks may miss important tasks that arrive during execution. In preemptive cases where the robot was physically manipulating the environment, additional overhead time is required to restore the robot’s pose and to continue grasping or other movement [Groth and Henrich, 2014]. Additionally, the robot must deal with failures; not only must the robot prioritize tasks, but it must decide (or have a plan for) what to do when the preempting task fails. Does the robot retry the failed task, move directly back to the preempted task, or drop into some kind of re-calibration or maintenance mode? Consider the Mars rover – if it runs into a rock or becomes stuck while navigating to a site where it has to perform chemical analysis, it should stop and get unstuck (or consult Earth-based humans for assistance), then return to navigation towards its earlier goal. If instead a piece of the rover’s chemical analysis fails due to hardware problems, it should probably stop all analyses until it can relay its problems and receive solutions from Earth.

Very limited literature exists on multi-tasking robots; much of the work focuses instead on robots that have many tasks to do very close together temporally, for example a UAV that searches for objects, targets an object, and releases a bomb. Groth and Henrich [2014] discusses a multi-tasking robot with nearby tasks (taking pictures of people and finding objects, or taking pictures of walls and greeting humans) with preemption; the robot stops taking pictures of walls if a human is in the way, for example.

In general, the MT-SR problem, regardless of the type of time window, can be approached heuristically as a bin packing problem, where each robot is a bin and each task is assigned to a robot that has available capacity and resources to perform that task. Other approaches to scheduling tasks are inspired by operat-
ing systems, such as shortest job first and priority scheduling; however, because
context switching is more time-intensive for robotics than it is for processors,
these methods may be highly suboptimal.

Lastly, as far as we know, no literature has explored the addition of stochas-
tic formulations with temporal constraints on the multi-task robot problem.
Probabilistic events, such as a delivery robot’s travel times due to traffic, would
affect the agent’s ability to perform single tasks, just as stochastic formulations
do in the ST-SR problem.

and Soft Constraints

Multi-task robots and multi-robot task problems remain sparsely explored
[Korsah et al., 2013, Gerkey and Matarić, 2004], even when additional temporal
constraints are not considered. This class of problems can be modeled as an
overlapping coalition formation problem [Chalkiadakis et al., 2010] combined
with a routing and scheduling problem. Standard coalition formation methods
produce either a super-coalition (with all the robots) or a set of non-overlapping
subsets of robots.

In cooperative games with overlapping coalitions, agents can do more than
one task at a time. This may lead robots to commit to the task assigned to more
than one coalition. Overlapping coalitions have been used to model collaborative
smartphone sensing in [Di et al., 2013]. In that work, smartphone users form
overlapping networks, and an incentive function rewards users’ contributions
to different tasks. Unfortunately, finding the optimal overlapping coalition is
NP-complete.

MT-MR-HC problems are comprised of the following subproblems: (1) as-
signing coalitions to tasks, (2) assigning different coalitions to the same robot
as long as no resource constraints are violated, and (3) assigning values to the
start and finishing times of tasks. Each of these subproblems is NP-hard.

Like MT-MR-HC, MT-MR-SC also lacks coverage in the MRTA literature.
The models used for MT-MR-HC can be extended to this class of problems,
the only difference being that temporal constraints are allowed to be violated.
This requires that one of the optimization objectives minimizes penalties from
violating the constraints.

6.7. Summarizing the Taxonomy

The coverage of our taxonomy is illustrated in Table 2. ST-SR-HC prob-
lems, both deterministic and stochastic, are by far the most commonly inves-
tigated. This is not surprising, given that problems in many application areas
can be modeled as single-robot and single-task problems with hard temporal
constraints. However, exogenous events do occur in robotic applications, hence
the lower prevalence of formal models for problems with soft constraints is sur-
prising.
### Table 2: Select papers from each category, with the percentage of papers that fall under each category (% Refs). The percentages are rough estimates from our references. Papers with a * symbol are not MRTA papers, but are included for completeness.

<table>
<thead>
<tr>
<th>Category</th>
<th>Deterministic % Refs</th>
<th>Stochastic % Refs</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>[Coltin and Veloso, 2014a]</td>
<td>[Ponda et al., 2012b]</td>
</tr>
<tr>
<td></td>
<td>[Gombolay et al., 2013]</td>
<td>[Pavone et al., 2011] *</td>
</tr>
<tr>
<td></td>
<td>[Barbulescu et al., 2010]</td>
<td>[Laporte et al., 1992] *</td>
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<td></td>
<td>[Ponda et al., 2010]</td>
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<td></td>
<td>[Jones et al., 2009]</td>
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<td></td>
<td>[Alighanbari et al., 2003]</td>
<td></td>
</tr>
<tr>
<td>ST-SR-SC</td>
<td>[Ta¸s et al., 2013] * 11</td>
<td>[Ta¸s et al., 2013] 2</td>
</tr>
<tr>
<td></td>
<td>[Hsu et al., 2007] *</td>
<td></td>
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<tr>
<td></td>
<td>[Ando and Taniguchi, 2006] *</td>
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</tr>
<tr>
<td></td>
<td>[Taillard et al., 1997] *</td>
<td></td>
</tr>
<tr>
<td>ST-MR-HC</td>
<td>Zhang and Parker, 2013b</td>
<td>13</td>
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<tr>
<td></td>
<td>Zhang et al., 2014</td>
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<tr>
<td></td>
<td>Zhang and Parker, 2013a</td>
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<td></td>
<td>Ramchurn et al., 2010b</td>
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<td></td>
<td>Parker and Tang, 2006</td>
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<tr>
<td>ST-MR-SC</td>
<td>[Amador et al., 2014] 3</td>
<td>0</td>
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<td></td>
<td>[Scerri et al., 2005]</td>
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<td></td>
<td>[Koes et al., 2005]</td>
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<tr>
<td>MT-SR-HC</td>
<td>Groth and Henrich, 2014</td>
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<tr>
<td>MT-SR-SC</td>
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</tr>
<tr>
<td>MT-MR-HC</td>
<td>Di et al., 2013 *</td>
<td>1</td>
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<tr>
<td>MT-MR-SC</td>
<td>0</td>
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</tr>
</tbody>
</table>

7. **Dynamic Task Release and Execution**

Execution of tasks in MRTA/TOC problems vary according to the dynamics considered. Dynamics may be due to faulty robots, changes in estimated cost of tasks due to uncertainties, changes in task definitions, online arrival of tasks, addition of robots to the team, and other changes made by external agents [Sariel-Talay et al., 2009]. While the execution aspect is outside of the task allocation scope, the planning-execution-replanning of tasks forms a planning loop that is usually addressed at once in dynamic domains. Here we consider dynamics caused by task arrival and during task execution separately.

Some dynamics are caused by the arrival of tasks over time without further knowledge of future tasks. Usually when a new task arrives at any given time there is already an existing solution for previously not yet performed but scheduled tasks. Thus, replanning occurs at task arrivals, while robots are executing previously assigned tasks [Cordeau and Laporte, 2007]. In [Nunes and Gini, 2015] both deterministic and dynamic task arrivals are considered, assuming the robots have perfect knowledge of the map where tasks appear. In contrast,
problems usually defined as online pickup and delivery problems or dial-a-ride include not only online arrival of tasks but other uncertain events, such as vehicle breakdowns and delays [Cordeau and Laporte, 2007]. Recent examples of online pickup and delivery consider transfers, in addition to the arrival of tasks with hard temporal constraints [Coltin and Veloso, 2014b,a, Bouros et al., 2011].

The dynamics that occur during plan execution [Usug and Sariel-Talay, 2011, Barbulescu et al., 2010, Ponda et al., 2010, Shah et al., 2009, Block et al., 2006] are very important for the practical use of robots, because execution can fail due to a host of reasons and replanning is essential to maintain some level of efficiency. In [Barbulescu et al., 2010] dynamics during execution are created by unexpected events and changes in costs and constraints; in [Ponda et al., 2010] dynamics are caused by breaks in communication links, which may cause conflicting assignments, as more than one robot could be assigned the same task. In [Usug and Sariel-Talay, 2011] temporary failures are considered, such as obstacles, which can be overcome by replanning.

8. Typical Solution Approaches

So far, we have proposed a taxonomy for MRTA/TOC problems; now we discuss the most popular solutions and how to map these to our taxonomy. Here we simply divide the methods into centralized vs. decentralized. Centralized methods are further separated into exact and approximate methods, while decentralized methods are grouped into distributed constraint-based and market-based according to nature of the proposed solutions.

8.1. Centralized Solutions

Centralized methods rely on a central controller that allocates tasks to robots. The autonomy of the robots in pure centralized methods is limited or non-existent, as they solely execute the dispatched orders and do not make decisions on what tasks to do.

MRTA/TOC is intractable for a non-trivial number of robots and tasks. Optimal centralized solutions are intractable because they need to evaluate a large number of candidate solutions in order to guarantee optimality. Thus, the focus of MRTA/TOC solutions is largely on approximation and heuristic solution methods. We discuss some of the common centralized exact and heuristic methods next.

8.1.1. Exact Solutions

Exact solutions are optimal, but their computation time is impractical for realistic robotics applications. The most naive way to search for such solutions is to exhaustively search for all possible allocations that do not violate the temporal constraints. This is, however, intractable, because an exhaustive search leads to worst-case $O(|K|!|A|)$ complexity for $|K|$ tasks and $|A|$ robots. We have to search through all the possible sequences of tasks and all possible allocations of tasks to robots, and in addition to all feasible assignments of times to tasks.
A more sophisticated method for MRTA/TOC problems is Branch-and-Bound (B&B) [Clausen, 1999]. B&B searches the state space of candidate solutions represented as a tree and uses upper and lower bounds of the optimal solution to prune the branches of the search tree that have costs higher than the computed lower bounds. Optimal solutions can be found using the B&B algorithm and its variants: Branch-and-Cut [Ropke et al., 2007, Bard et al., 2002], Branch-and-Price [Korsah et al., 2012, Feillet, 2010, Dohn et al., 2009, Barnhart et al., 1998], and Branch-Price-and-Cut [Barnhart et al., 2000]. Branch-Price-and-Cut is becoming more popular in VRPTW [Bettinelli et al., 2011, Archetti et al., 2011, Desaulniers, 2010, Ropke and Cordeau, 2009], but, as far as we know, this technique has not been used in MRTA/TOC problems.

For many of these exact methods problems are solved using tools such as CPLEX [ILOG, 2006], Gurobi [Gurobi Optimization, 2014], ABACUS [Jünger and Thienel, 2000], lp solve [Berkelaar et al., 2004] or other tools to build and solve the underlyng MILP formulations. MILP-based formulations and solutions have been predominantly used in ST-SR-HC problems [Korsah et al., 2012, Alighanbari et al., 2003], but these models have also been used in other parts of the taxonomy (e.g. ST-MR-HC [Ramchurn et al., 2010b, Koes et al., 2005]).

8.1.2. Approximate and Heuristic Solutions

To reduce computation time, MILP-based heuristics are used to find approximate partial allocations while searching the state-space tree. Such approaches have been used to address problems in our taxonomy (e.g. ST-SR-HC [Gombolay et al., 2013, Korsah et al., 2012]). As far as we know, these methods do not provide any theoretical guarantees, but in some cases (e.g. [Gombolay et al., 2013]) they experimentally achieve results that are only 10% away from the optimal value (makespan).

Another way to gain computational efficiency is to use metaheuristic approaches. Metaheuristics are algorithmic templates that approximately solve hard combinatorial optimization problems. Unlike other combinatorial optimization algorithms, metaheuristics may allow lower quality solutions in the search process to escape local optima, and often embed off-the-shelf heuristics to solve the problem [Bräysy and Gendreau, 2005b].

Metaheuristic approaches to VRPTW, TOPTW and related routing and scheduling problems have been shown to outperform many other methods (e.g. construction heuristics and local search) for standard benchmarks [Bräysy and Gendreau, 2005b, Hu and Lim, 2014]. Recent trends in the metaheuristic literature seek to reduce the computation time and improve the solution quality by using parallelization and hybridization of different heuristics and exact techniques. However, metaheuristic parameters remain hard to tune [Birattari, 2009, Bräysy and Gendreau, 2005b].

8.2. Decentralized Solutions

Decentralized approaches vary widely; a detailed categorization is outside the scope of this paper and we refer the reader to [Dias et al., 2006, Pentico, 2007,
Koenig et al., 2010] for more thorough taxonomies on MRTA methods. Here we focus on distributed constraint optimization and market- and negotiation-based algorithms since these have received a great deal of attention in the MRTA community.

8.2.1. Distributed Constraint (DCOP)-Based Methods

MRTA/TOC problems can be modeled as a Distributed Constraint Optimization Problem (DCOP) [Maheswaran et al., 2004] and solved using DCOP methods. Solving DCOP exactly is NP-hard and impractical even for unconstrained MRTA problems [Junges and Bazzan, 2008]. Thus, approximate methods such as Max-Sum [Farinelli et al., 2008] have been used for task allocation in sensor networks and in RoboCup Rescue [Ramchurn et al., 2010a].

Ramchurn et al. [2010a] proposed the Fast Max-Sum algorithm, which was shown to be robust in situations where the number of tasks is dynamic; the approach reduced the computation time, number and size of messages sent compared to Max-Sum. The work in [Macarthur et al., 2011] improved upon the DCOP solution proposed in [Ramchurn et al., 2010a] by using online domain pruning and branch-and-bound. Their solution uses less computation overhead in a dynamic environment.

Another method is the LA-DCOP [Farinelli et al., 2006] algorithm, which is an approximation that uses token passing [Xu et al., 2005] as follows: when an agent perceives a task, it creates a token to represent it. It can decide to do the task or pass the token to a randomly chosen agent. This tends to guide the search quickly towards a greedy solution, which is reasonable for ST-SR-HC problems.

In [Ferreira et al., 2008] LA-DCOP and Swarm-GAP are compared in RoboCup settings. In Swarm-GAP an agent chooses a task according to a probability that depends on the stimulus generated by the task and the agent’s threshold. Results show that both DCOP approaches behave similarly, and both perform better than a greedy task allocation. Their approach works for ST-MR-SC problems, where agents are allowed to arrive late to tasks. Recently, to facilitate comparing the performance of DCOP algorithms, RMASBench, a system that provides a library of state-of-the-art solvers for DCOP and for comparing them, has been created in [Kleiner et al., 2013].

8.2.2. Market and Negotiation-Based Methods

Among the decentralized algorithms, sequential auction- and negotiation-based algorithms [e.g. Nunes and Gini, 2015, Sariel-Talay et al., 2009, Ponda et al., 2010] are more prevalent than other methods. Sequential auction algorithms produce solutions that are two away from optimal in the worst-case in both single-item [Lagoudakis et al., 2004] and multi-item auctions [Choi et al., 2009]. This, together with the ease of implementation and extension to dynamic scenarios and robust execution [Nanjanath and Gini, 2010] makes sequential auctions an attractive solution. However, the greedy nature of sequential auctions and the complex structure of most MRTA/TOC problems cause the addition of temporal constraints to auction algorithms to produce suboptimal solutions.
Temporal modeling and balancing between temporal- and distance-based objectives can help auctions perform better [Nunes and Gini, 2015, Ponda et al., 2010]. In [Amador et al., 2014] Fisher markets are used in a decentralized online solution to dynamic ST-MR-SC problems were agents are considered buyers and tasks are the goods.

Auctions distribute the computation to individual agents but require communication to share bids and results. To reduce the need for communication, several approaches use consensus algorithms [Zavlanos et al., 2008, Choi et al., 2009, Ponda et al., 2010], where each agent determines independently which tasks it should do. An equilibrium is reached by iteratively sharing information with neighbors and re-allocating tasks if needed. [Godoy and Gini, 2012] extended the Consensus Based Bundle Algorithm (CBBA) [Choi et al., 2009] to optimize the number of completed tasks for tasks with temporal constraints in ST-SR-HC problems. A different method, called emergent task allocation [Atay and Bayazit, 2003], distributes the computation of task allocation for a surveillance task to individual robots, by sharing intentions and directives with 1-hop away neighbors. The method has been shown to converge to the optimal solution as the number of iterations of information sharing increases.

Despite the development of many decentralized methods for MRTA/TOC problems, very limited work offers theoretical analysis of the quality of these solutions. There is a need for theoretical performance bounds for both centralized and decentralized heuristics for the MRTA/TOC problem.

There are other decentralized approaches to task allocation that are not market- or DCOP-based. For instance, [Chapman et al., 2010] formulated ST-MR MRTA as a stochastic game and used overlapping potential games to approximate an optimal solution. Their approach is robust to restricted agent communication and observation range.

Swarm-based approaches have been proposed for various tasks, such as foraging, where robots need to find food and bring it to the nest [Lerman et al., 2006, Brutschy et al., 2014] or where swarms of robots are allocated different monitoring tasks without any communication among the robots [Berman et al., 2009]. Swarm methods often work well but do not have theoretical guarantees.

9. Summary, Open Issues, Direction for Future Research

Problems that consider temporal and ordering constraints relate to many well studied problems, such as vehicle routing, job-shop scheduling, and multirobot task allocation. A large portion of the literature in MRTA/TOC focuses on ST-SR-HC problems, some address the soft constraint version of this class of problem; however, the literature is sparser for other classes of problems that consider multi-task robots and multi-robot tasks.

9.1. Summary

We surveyed the multi-robot task allocation literature related to problems where tasks have constraints on where, when and possibly the order in which
they have to be performed. We built on a previous taxonomy and added a
classification axis that separates the literature according to the hard versus soft
nature of the constraints. Where appropriate, we gave a generic mathematical
formulation of problems both in deterministic and stochastic cases, and
offered an account of some common execution dynamics. We briefly discussed
the methods applied to the problems in our taxonomy, and split solutions into
centralized and decentralized approaches. In addition, our work drew parallels between multi-robot task allocation with temporal and ordering constraints
with other areas of research, and throughout the paper we discussed models and
solutions coming from these areas. Lastly, our work discussed areas that are still
sparsely covered, and provided directions for future research in the area.

9.2. Open Issues and Future Research

There are several open issues that need to be addressed, which we did
not exhaustively address here. Progress in the following topics would greatly
advance research in MRTA/TOC: (1) models and algorithms for stochastic
MRTA/TOC problems, (2) study of theoretical guarantees of approximate sol-
uations, (3) richer and more complex temporal models with provably good and
efficient algorithms, (4) models and algorithms for multi-task robot task alloca-
tion problems, (5) studies on the effects of time scales and time sensitivity in
MRTA/TOC problems and (6) the development of a research platform to make
software and data available to researchers.

Research in stochastic MRTA/TOC problems is still very sparse. The de-
velopment of MRTA methods that take advantage of simulation and stochastic
models to better plan under uncertainty is an endeavor worth pursuing because
robots often operate in uncertain environments. Important research questions
can be asked here; for example, in an uncertain environment is it more benefi-
cial to build a complex model that incorporates uncertainty, or is it enough to
build less well-informed plans and replan as often as needed to quickly react to
unexpected events?

There is also a need for work on theoretical guarantees for many heuristic
schedulers developed for MRTA/TOC problems. The NP-complete nature of the
problem and the need for relatively fast planners has generated many heuristics.
However, such heuristics typically lack performance guarantees, which can be
crucial for safety critical systems, to ensure that robots work effectively even in
the worst possible scenarios.

More work needs to be done to address more complex temporal constraint
types, such as disjunctive temporal models. The literature could also benefit
from work that combines soft and hard time windows, and precedence with
simultaneity constraints. A mix of these constraints might produce more ex-
pressive models for a larger set of real-world problems.

The challenge of allocating tasks to multi-task robots, which are robots that
can perform more than one task at a time, remains open. As we can imagine this
is difficult for many existing robots because they lack the necessary actuators.
The lack of literature might also be due to the lack of practical applications
of multi-task robots. We are not aware of any practical problems that strictly
requires robots to perform multiple tasks concurrently; one example of such
problems could be in military domains where a drone robot could be required
to strike a target while at the same tracking other targets in nearby areas.

Another interesting, yet not so explored topic, regards time scales and sensi-
tivity. Robots that move and navigate in an environment can be on a short time
scale, as in exploring a building in hours, or a large time scale of exploring a
planet for years. Even for a single robot, tasks can have varied time sensitivities;
some tasks may have short hard time constraints, whereas others may have long
time horizons with soft time windows. For instance, the Mars rover Curiosity
has periodic tasks that occur every day for years (data upload), constantly run-
ning tasks (temperature regulation), and sporadic tasks (chemical analysis of
collected material and drilling). Each of these tasks have different time sensi-
tivities; for example, data upload needs to occur when the receiving orbiters are
within view of the rover. Temperature regulation requires constant vigilance,
and drilling can be postponed, but needs to occur when the rover is within reach
of the material. Chemical analysis in the rover’s internal chambers can occur
regardless of location. This single robot has tasks with hard time windows, soft
time windows, varied scheduling horizons, and varied sensitivities. Considering
tasks in terms of time scales and task sensitivities in the same robotic system
thus holds value for any researcher interested in real world problems.

Lastly, we are concerned with the public availability of research data and
methods. We advocate for a computational infrastructure for MRTA problems
(in general, or in particular problems with temporal and ordering constraints).
A tool identical to the Computational Infrastructure for Operations Research
COIN-OR [2015]) could greatly benefit MRTA researchers. COIN-OR is an
open source software project in which many operations research algorithms are
implemented and maintained by scholars in the area. That in combination with
datasets would help researchers verify their results on publicly available data
and methods, allowing for richer comparisons among methods.

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